

# New RSA Vulnerabilities Using Lattice Reduction Methods

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## Abstract

In this dissertation thesis we study today's most popular and widely used public-key cryptosystem: the RSA-cryptosystem, which was proposed in 1978 by Rivest, Shamir and Adleman. We show that some choices of the RSA-parameters lead to polynomial time attacks on the cryptosystem.

Let us have a closer look at the parameter generation process in RSA: One chooses two large prime numbers  $p$  and  $q$  and computes their product  $N = pq$ . The so-called RSA modulus  $N$  is public, whereas the factorization of  $N$  is kept secret. Furthermore, one chooses a key-pair  $(e, d)$  satisfying  $ed = 1 \pmod{(p-1)(q-1)}$ . The parameter  $e$  is public and the parameter  $d$  is secret.

One can easily compute the secret key  $d$  from the public information  $(N, e)$  if the factorization of  $N$  is known. Therefore, an attacker can try to compute the factorization of  $N$ . But up to now no algorithm is known that factors  $N$  in time polynomial in the bit-size of  $N$ .

In this work, we show that an attacker can determine the factorization of  $N$  in polynomial time, provided that  $e$  is of a special form or that the attacker gets into possession of a fraction of the secret key bits. The main method that we use in order to achieve our results is a method for finding small solutions of modular polynomial equations, which was proposed in 1996 by Coppersmith. Our main results in this dissertation thesis are:

- Generalization of Coppersmith's method for univariate polynomials. We also propose an approach for constructing optimal lattice bases, which are used in Coppersmith's method.
- Polynomial time factorization of  $N$  given  $(N, e)$ , provided that the corresponding  $d$  is of the form  $d = \frac{d_1}{d_2} \pmod{(p-1)(q-1)}$  for small  $d_1, d_2$ .

This result leads to the cryptanalysis of an RSA-variant that was proposed in 2001.

- Polynomial time factorization of  $N$ , provided that the parameter  $d_p = d \pmod{p-1}$  is small and that  $q \leq N^{0.382}$ . Small values of  $d_p$  are often used in practice since they speed up the decryption process.
- Several polynomial time attacks on RSA, provided that a fraction of the bits of  $d$  or of  $d_p$  is known. These attacks are also generalized to RSA moduli of the form  $N = p^r q$  for  $r > 1$ .